

A New Approach To Efficient Revocable Attribute-Based Anonymous Credentials

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Outline

1. Introduction
2. Novel ABC Paradigm [HS14]
3. Novel Revocation Approach
 - Security Model
 - Constructions
4. Conclusion

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Attribute-Based Anonymous Credentials (ABCs)

Organization

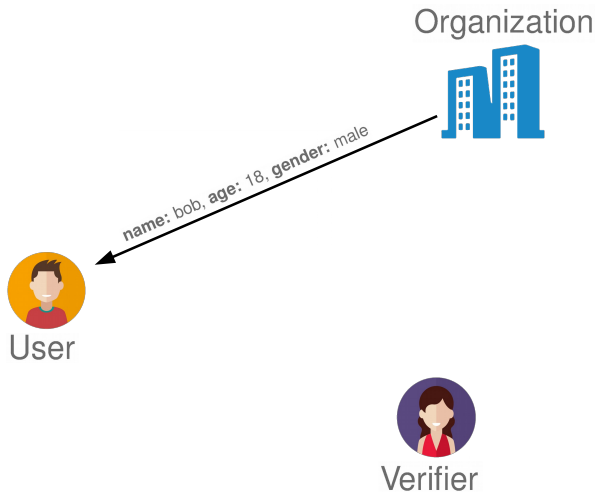


User

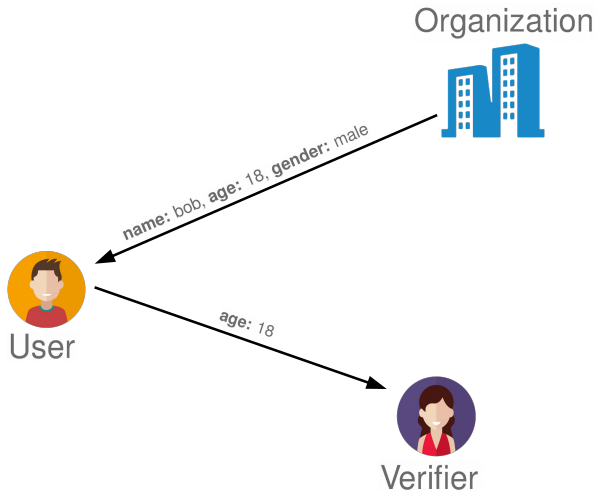


Verifier

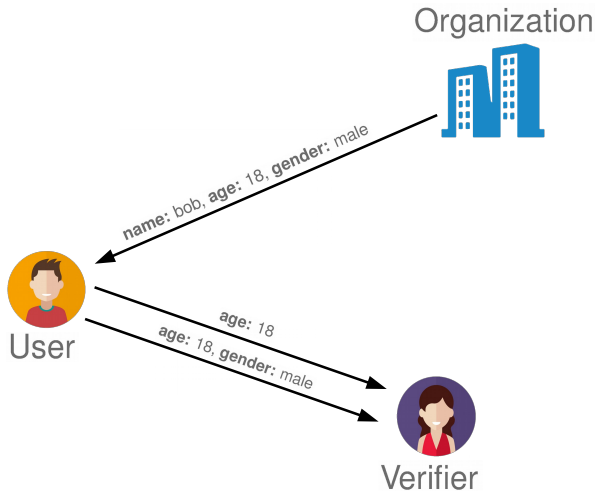
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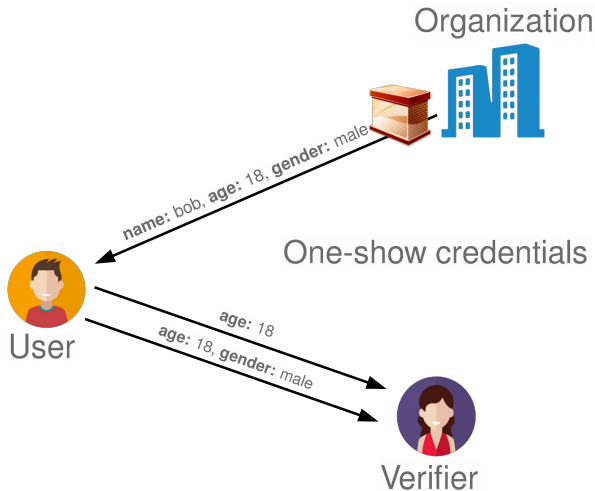
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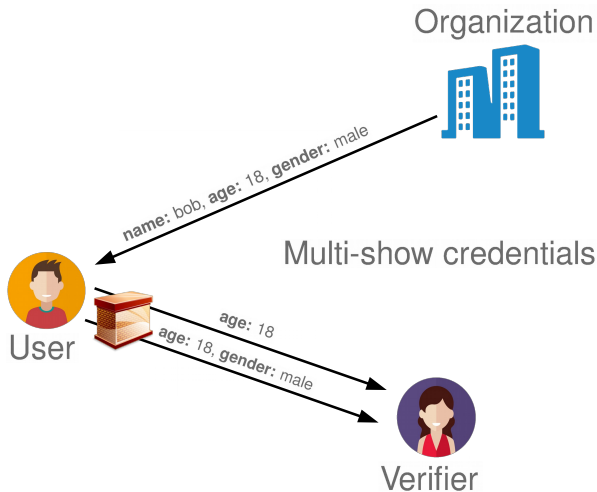
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- Most prominent: IBM's Idemix, Microsoft's U-Prove

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Novel approach to multi-show ABCs [HS14]

- **Structure Preserving Signatures on Equivalence Classes**
 - (SPS-EQ)
- Interesting properties
 - $O(1)$ size of creds and $O(1)$ communication
 - No PoK for unrevealed attributes
 - Only single $O(1)$ PoK for freshness

Motivation

Revocation: **important** feature in practice

Question

Revocation mechanism

- Following similar principles
- Preserving the nice asymptotic properties

Contribution

Accumulator-based **blacklist revocation** for [HS14]

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Revocation **typically** considered as an **add-on**

- Extend model of [HS14]
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Comparison of both approaches

Preliminaries

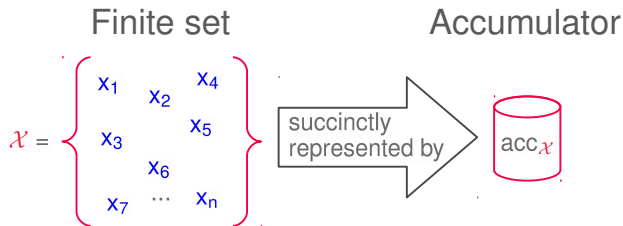
Asymmetric bilinear map (pairing)

- $e : \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{G}_T$, where $|\mathbb{G}_1| = |\mathbb{G}_2| = |\mathbb{G}_T| = p$
- $e(aP, b\hat{P}) = e(P, \hat{P})^{ab}$ (Bilinearity)
- $e(P, \hat{P}) \neq 1_{\mathbb{G}_T}$ (Non-degeneracy)
- $e(\cdot, \cdot)$ efficiently computable (Efficiency)

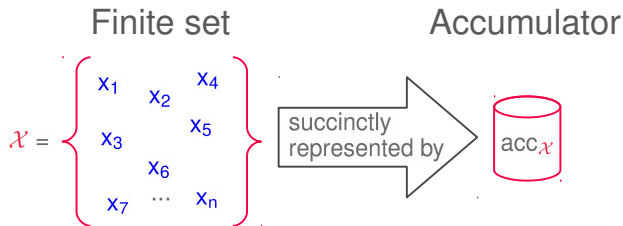
SXDH setting

⇒ DDH assumed to hold in \mathbb{G}_1 and \mathbb{G}_2

Cryptographic Accumulators



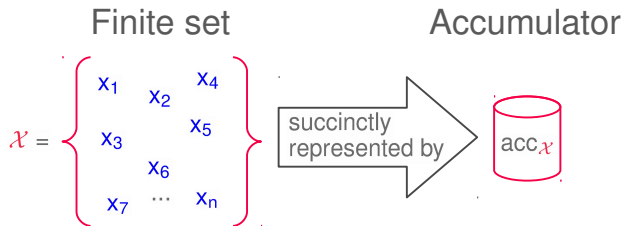
Cryptographic Accumulators



Witnesses wit_x certifying membership of x in acc_X

- Efficiently computable $\forall x \in \mathcal{X}$
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\Rightarrow **Collision freeness**

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Structure Preserving Signatures [AFG⁺10]

Sign group element vectors

⇒ Sigs and PKs only consist of group elements

Verification solely via

- Pairing-product equations
- Group membership tests

Signing Equivalence Classes [HS14]

Partition \mathbb{G}_i^ℓ into projective equivalence classes

$$M \in \mathbb{G}_i^\ell \sim_{\mathcal{R}} N \in \mathbb{G}_i^\ell \Leftrightarrow \exists k \in \mathbb{Z}_p^* : N = k \cdot M$$

SPS-EQ

- Given σ on M
 - Publicly derive σ' on $M' \in [M]_{\mathcal{R}}$
- \Rightarrow IND of classes under DDH in \mathbb{G}_i

Signing Equivalence Classes [HS14] II

Security properties

- Correctness
- EUF-CMA security (w.r.t. equivalence classes)
- Perfect adaption of signatures
 - (M', σ') obtained by re-randomizing (M, σ)
 - Indistinguishable from fresh signature on M'

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Credential

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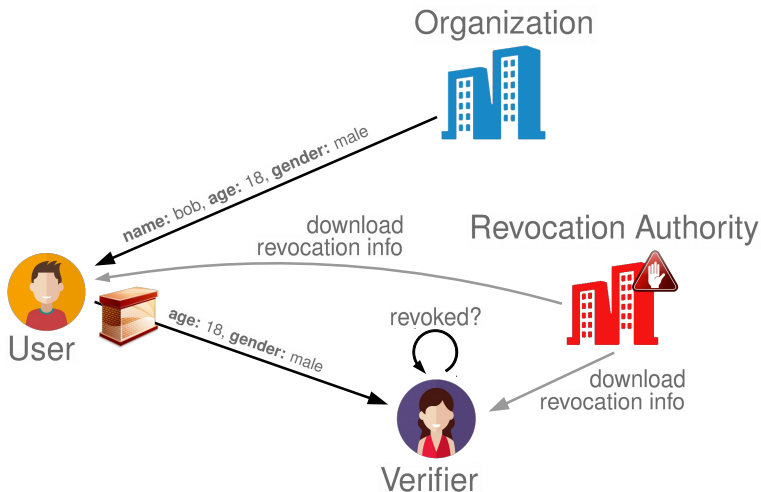
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⇒ Adversary:

- Can corrupt users, obtain secret keys of users, revoke users
- Acts as dishonest user

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Novel Revocation Approach

Idea:

- Choose accumulator s.t. compatible with re-randomization

Construction (sketch):

- Incorporate **nym** as credential component
- Accumulator contains all revoked *nyms*

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 - Consistently randomize accumulator/witnesses
 - Plug-in randomized **nym** into verification relation

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 - Consistently randomize accumulator/witnesses
 - Plug-in randomized **nym** into verification relation
- + Simple $O(1)$ PoK for technical reasons

Security

Unforgeability

- \approx Unforgeability of underlying ABC
- + Case for collision freeness of accu
- + Some technicalities regarding extraction
 - 3 additional DLOG proofs

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Anonymity contd'

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DDH-like assumptions

- 1 holds in GGM
- 1 follows from (R, S, T, f) -DDH [Boy08]

U-Prove Based Revocation Approach

Follow **classical** revocation approach

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Follow **classical** revocation approach

- Adapt U-Prove revocation [ACN13, NP14]

Construction (sketch):

- Incorporate **nym** in credentials
- Accumulator contains revoked *nyms*
- PoK of non-membership witness and **nym**
- PoK that **nym** coincides with **nym** in credential

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Anonymity

- Components depending on the challenge bit
 - Indistinguishable from random
 - Shown under DDH in \mathbb{G}_1
 - ... and perfect adaption of signatures

Comparison

Comparison based on [UW14]

- BN implementation on ARM-Cortex-M0+

Obtain:

- Novel Approach: +15 G_1 equivalents
- Classic Approach: +14 G_1 equivalents

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- BN implementation on ARM-Cortex-M0+

Obtain:

- Novel Approach: +15 G_1 equivalents
- Classic Approach: +14 G_1 equivalents

Show:

- Novel Approach: +20 G_1 equivalents
- Classic Approach: +33 G_1 equivalents

⇒ worst case (best case: even $2 \times$ faster)

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New conceptually simple approach

- Easy to comprehend
 - Easy to implement
- ⇒ New direction in revocation for ABC systems

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Performance

- Both approaches practically efficient
- Novel approach yields more efficient showings

Thank you.

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Supported by **prisma cloud**

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